

# Topics in Experimental and Behavioral Economics

## Quantal Response - A Noisy Equilibrium

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### Introduction

1. Quantal Response Equilibrium
2. Logit equilibrium
3. Some experiments
4. Logit equilibrium in continuous strategies
5. Examples and applications

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## Quantal Response Equilibrium

McKelvey & Palfrey GEB95

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- Consider a **Normal Form Game**  $\Gamma = \{N, S, u\}$
- $n$  players:  $N = \{1, \dots, n\}$   
strategy set  $S_i = \{s_{i1}, \dots, s_{iJ_i}\}$ ,  $J_i$  pure strategies of player  $i$   
payoff function  $u_i : S \rightarrow \mathcal{R}$ , where  $S = \prod_{i \in N} S_i$
- $J_i$ -dim simplex  $\Delta_i = \{p_i = (p_{i1}, \dots, p_{iJ_i}) : \sum p_i = 1, p_{ij} \geq 0\}$ ;  
 $p_i$  is a mixed strategy for player  $i$
- $\Delta = \prod_{i \in N} \Delta_i$  space of mixed strategies;  
 $p = (p_1, \dots, p_n)$  vector of mixed strategies
- payoff  $u_i(p) = \sum_{s \in S} p(s) u_i(s)$ , where  $p(s) = \prod_{i \in N} p_i(s_i)$   
(expected payoff when players play mixed strategies  $p$ )

## QRE - Definitions

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- **Definition of Nash Equilibrium:**  
A vector  $p = (p_1, \dots, p_n) \in \Delta$  of (mixed) strategies is a *Nash equilibrium* if for all  $i \in N$  and all  $p'_i \in \Delta_i$ ,  $u_i(p'_i, p_{-i}) \leq u_i(p)$

## QRE - Definitions

- Random utility - Utility with noise
  - space of possible payoffs of player  $i$ :  
 $X_i = R^{J_i}$ ;  $X = \prod_{i=1}^n X_i$
  - Define  $\bar{u} : \Delta \rightarrow X$  as

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$$\bar{u}(p) = (\bar{u}_1(p), \dots, \bar{u}_n(p)) \text{ where } \bar{u}_{ij}(p) = u_i(s_{ij}, p_{-i})$$

- player  $i$ 's utility is subject to random error:

$$\hat{u}_{ij}(p) = \bar{u}_{ij}(p) + \epsilon_{ij}$$

where  $\epsilon_i = (\epsilon_{i1}, \dots, \epsilon_{iJ_i})$  player  $i$ 's error vector, with  $E[\epsilon_i] = 0$  and density function  $f_i(\epsilon_i)$ .

## QRE - Definitions

- **Behavioral assumption:**
  - each player  $i$  selects an action  $j$  such that  $\hat{u}_{ij} \geq \hat{u}_{ik}$  for all  $k = 1, \dots, J_i$
- Define  $ij$ -response set  $R_{ij} \subseteq R^{J_i}$  as

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$$R_{ij}(\bar{u}_i) = \{\epsilon_i \in R^{J_i} \mid \bar{u}_{ij} + \epsilon_{ij} \geq \bar{u}_{ik} + \epsilon_{ik} \text{ for all } k = 1, \dots, J_i\}$$

(region of errors that will lead  $i$  to choose  $j$ , for given  $p$ )

- **Statistical response function** gives the probability that player  $i$  will choose action  $j$ , given  $\bar{u}$  and  $f$

$$\sigma_{ij}(\bar{u}_i) = \int_{R_{ij}(\bar{u}_i)} f(\epsilon) d\epsilon$$

## QRE - Definitions

- Definition of **Quantal Response Equilibrium**:

For any admissible  $f = (f_1, \dots, f_n)$  and game  $\Gamma = (N, S, u)$  a *quantal response equilibrium* (QRE) is any vector  $\pi \in \Delta$  such that for all  $i \in N$  and all  $1 \leq j \leq J_i$ ,  $\pi_{ij} = \sigma_{ij}(\bar{u}_i(\pi))$ .

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$\sigma_i : R^{J_i} \rightarrow \Delta^{J_i}$  is the *statistical reaction function* (or *quantal response function*) of player  $i$ .

- QRE is a statistical version of NE, where each agent's expected utility is subject to a random error, e.g. due to calculation error.

## QRE - Properties

- **Properties of the statistical reaction function:**

1.  $\sigma \in \Delta$  is non-empty.
2.  $\sigma_i$  is continuous on  $R^{J_i}$ .
3.  $\sigma_{ij}$  monotonically increasing in  $\bar{u}_{ij}$ .
4.  $\bar{u}_{ij} > \bar{u}_{ik} \Rightarrow \sigma_{ij}(\bar{u}) > \sigma_{ik}(\bar{u})$ .

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- 1. and 2. imply **existence of QRE**, for any admissible  $f$ .
- 3. and 4. imply 'better actions are more likely to be chosen'.

### Logit Equilibrium

- A simple parametric class of QRE - **Logit Equilibrium**:  
If errors are independent with extreme value distribution, then for payoff  $x_i \in \mathcal{R}^{J_i}$  the **logistic** quantal response function is given by

$$\sigma_{ij}(x_i) = \frac{e^{\lambda x_{ij}}}{\sum_{k=1}^{J_i} e^{\lambda x_{ik}}}$$

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- QRE or **Logit Equilibrium** requires

$$\pi_{ij} = \frac{e^{\lambda x_{ij}}}{\sum_{k=1}^{J_i} e^{\lambda x_{ik}}}, \quad \text{where } x_{ij} = \bar{u}_{ij}(\pi)$$

parameter  $\lambda$  inversely related to level of error:

$\lambda = 0$ : all error; all probabilities equal  $1/J_i$

$\lambda = +\infty$ : no error; Logit Eq. approaches NE

### Logit Equilibrium

- **Properties of the Logit Equilibrium**:
- Define the **Logit Equilibrium correspondence**  
 $\pi^* : \mathcal{R}^+ \rightarrow 2^\Delta$  by

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$$\pi^*(\lambda) = \left\{ \pi \in \Delta : \pi_{ij} = \frac{e^{\lambda \bar{u}_{ij}(\pi)}}{\sum_{k=1}^{J_i} e^{\lambda \bar{u}_{ik}(\pi)}} \quad \forall i, j \right\}$$

- Note: For  $\lambda = 0$

$$\pi^*(0) = \left\{ \pi \in \Delta : \pi_{ij} = \frac{e^0}{\sum_{k=1}^{J_i} e^0} \quad \forall i, j \right\} = (1/J_1, 1/J_2, \dots, 1/J_n)$$

## Logit Equilibrium

- **Theorem:** Let  $\sigma$  be the logistic quantal response function, and  $\{\lambda_1, \lambda_2, \dots\}$  be a sequence s.th.  $\lim_{t \rightarrow \infty} \lambda_t = \infty$ . Let  $\{p_1, p_2, \dots\}$  be a corresponding sequence with  $p_t \in \pi^*(\lambda_t) \forall t$ , s.th.  $\lim_{t \rightarrow \infty} p_t = p^*$ , then  $p^*$  is a Nash equilibrium.

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- **Theorem:** For almost all games  $\Gamma = (N, S, u)$ 
  1.  $\pi^*(\lambda)$  is odd for almost all  $\lambda$
  2.  $\pi^*$  is upper hemicontinuous
  3. The graph of  $\pi^*$  contains a unique branch from the centroid for  $\lambda = 0$  to a unique NE as  $\lambda$  goes to  $+\infty$

## Logit Equilibrium and experiments

- Some first laboratory testing of the Logit Equilibrium for two-person games with a unique Nash Equilibrium and without payoffs Pareto preferred to the NE

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- Idea: Calculate maximum likelihood estimate of  $\lambda$  and check how well it fits the data.
- Analyzed experiments span more than 30 years  $\rightarrow$  all payoffs converted to real 1982 dollars

### Logit Equilibrium and experiments

- Laboratory Experiment I (Lieberman, 1960) - Two person zero sum game:

	$B_1$	$B_2$	$B_3$
$A_1$	15,-15	0,0	-2,2
$A_2$	0,0	-15,15	-1,1
$A_3$	1,-1	2,-2	0,0

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- unique Nash equilibrium ( $A_3, B_3$ )
- each subject in 200 plays with single opponent

### Logit Equilibrium and experiments

- Lieberman zero sum game:

	$B_1$	$B_2$	$B_3$
$A_1$	15,-15	0,0	-2,2
$A_2$	0,0	-15,15	-1,1
$A_3$	1,-1	2,-2	0,0

TABLE III  
DATA AND ESTIMATES FOR LIEBERMAN (1960) EXPERIMENTS:  $N = 300$  FOR EACH EXPERIENCE LEVEL

Periods	Actual Data				Predicted			$\lambda$	$-\hat{U}^*$
	$A_1$	$A_2$	$B_2$	$B_3$	$\hat{A}_1$	$\hat{A}_2$	$\hat{B}_3$		
1-10	0.260	0.720	0.300	0.667	0.277	0.696	0.176	212.0	
11-20	0.167	0.806	0.227	0.760	0.196	0.781	0.252	177.0	
21-30	0.113	0.880	0.160	0.833	0.138	0.838	0.329	134.3	
31-40	0.093	0.887	0.120	0.853	0.106	0.869	0.390	134.4	
41-50	0.060	0.907	0.073	0.907	0.066	0.906	0.500	109.5	
51-60	0.060	0.873	0.120	0.860	0.087	0.886	0.435	144.7	
61-70	0.060	0.853	0.113	0.867	0.083	0.890	0.448	152.7	
71-80	0.060	0.907	0.047	0.933	0.054	0.916	0.547	98.9	
81-90	0.047	0.893	0.067	0.920	0.056	0.915	0.542	112.3	
91-100	0.027	0.920	0.080	0.907	0.053	0.918	0.553	105.6	
101-120	0.053	0.907	0.047	0.933	0.051	0.920	0.564	99.5	
111-120	0.027	0.920	0.047	0.933	0.037	0.932	0.635	94.2	
121-130	0.040	0.927	0.040	0.920	0.040	0.929	0.616	97.1	
131-140	0.033	0.927	0.047	0.953	0.040	0.929	0.616	80.2	
141-150	0.053	0.913	0.060	0.900	0.056	0.915	0.542	112.3	
151-160	0.053	0.900	0.053	0.920	0.052	0.919	0.558	109.3	
161-170	0.027	0.946	0.060	0.927	0.045	0.925	0.592	83.4	
171-180	0.053	0.900	0.033	0.927	0.042	0.927	0.604	107.1	
181-190	0.027	0.933	0.020	0.973	0.023	0.946	0.737	67.0	
191-200	0.040	0.920	0.047	0.933	0.044	0.926	0.598	93.7	

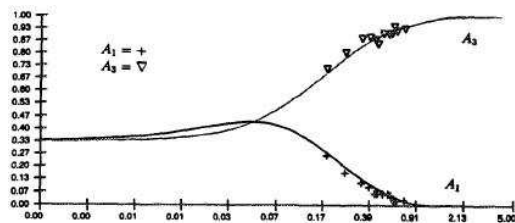
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- data broken down to 20 experience levels of 10 periods QRE estimates for each level
- during early rounds player 1 overplays  $A_1$  and player 2 overplays  $B_2$ , relative to NE; both predicted by QRE (see fig. next slide)

### Logit Equilibrium and experiments

- Lieberman zero sum game:

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- QRE prediction and actual data

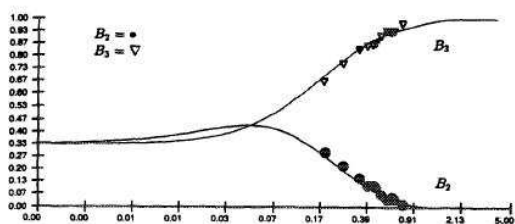


FIG. 3. QRE as a function of  $A$  for Lieberman experiment.

### Logit Equilibrium and experiments

- Laboratory Experiment II (O'Neill, 1987) - Zero sum game

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	$B_1$	$B_2$	$B_3$	$B_4$
$A_1$	5,-5	-5,5	-5,5	-5,5
$A_2$	-5,5	-5,5	5,-5	5,-5
$A_3$	-5,5	5,-5	-5,5	5,-5
$A_4$	-5,5	5,-5	5,-5	-5,5

- unique Nash equilibrium in mixed strategies:  
(0.4, 0.2, 0.2, 0.2) both players
- each subject in 105 plays

### Logit Equilibrium and experiments

- O'Neill zero sum game:

TABLE IV  
DATA AND ESTIMATES FOR O'NEILL

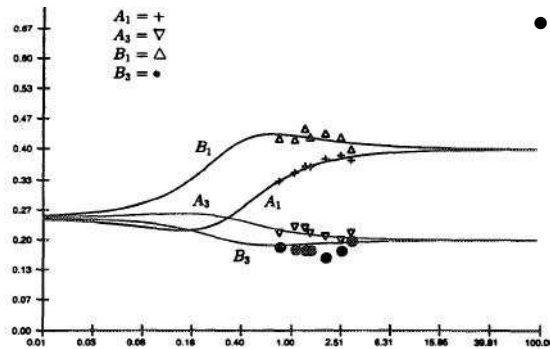
	Number	Frequency	Rand	NE	QRE
$A_1$	949	0.362	0.250	0.400	0.360
$A_2$	579	0.221	0.250	0.200	0.213
$A_3$	565	0.215	0.250	0.200	0.213
$A_4$	532	0.203	0.250	0.200	0.213
$B_1$	1119	0.426	0.250	0.400	0.426
$B_2$	592	0.226	0.250	0.200	0.191
$B_3$	470	0.179	0.250	0.200	0.191
$B_4$	444	0.169	0.250	0.200	0.191
$\lambda$			0	$\infty$	1.313
$-\lambda^*$			7278	7016	7004

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- data broken down to 7 experience levels of 15 periods
- data show:  $A_1$  is 'underplayed';  $B_1$  is 'overplayed'
- QRE predicts:  $A_1 < B_1$  is 'overplayed'
- QRE does significantly better than Nash and Random model; LR test

### Logit Equilibrium and experiments

- O'Neill zero sum game:



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- QRE prediction:  
for intermediate  $\lambda$ :  $B_1$  overplayed,  $A_1$  underplayed  
for all  $\lambda$ :  $B_1 > A_1$

FIG. 4. QRE as a function of  $\lambda$  for O'Neill experiment.

### Logit Equilibrium and experiments

- Laboratory Experiment IV (Ochs, 1995) - (Non-)zero-sum games

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Game 1			Game 2			Game 3		
	$B_1$	$B_2$		$B_1$	$B_2$		$B_1$	$B_2$
$A_1$	1,0	0,1	$A_1$	9,0	0,1	$A_1$	4,0	0,1
$A_2$	0,1	1,0	$A_2$	0,1	1,0	$A_2$	0,1	1,0

- Only difference is payoff to player 1 if  $(A_1, B_1)$  is played.
- In the unique Nash equilibrium:  
 player 1 mixes always with 0.5  
 player 2 plays  $B_1$  with probability 0.5, 0.1, and 0.2,  
 respectively.

### Logit Equilibrium and experiments

- Ochs 95 - Non-zero-sum game:  
 52 plays, 4 experience levels of 16 periods

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TABLE IX  
DATA AND ESTIMATES FOR OCHS, GAME 2

Period	$n$	Actual		Predicted		$\lambda$	$-\hat{u}^*$		
		$A_1$	$B_1$	$\hat{A}_1$	$\hat{B}_1$		QRE	Nash	Rand
1-16	128	0.541	0.326	0.645	0.347	1.951	1721	1938	1774
17-32	128	0.649	0.228	0.645	0.228	3.763	1517	1664	1774
33-48	128	0.578	0.250	0.648	0.241	3.475	1605	1725	1774
48-52	64	0.626	0.200	0.636	0.197	4.638	743	792	887
All	448	0.595	0.258	0.649	0.254	3.241	5612	6119	6210

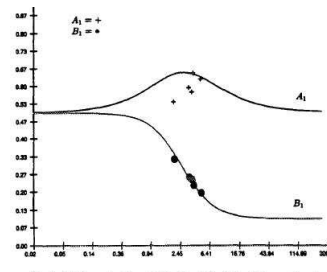


FIG. 6. QRE as a function of  $\lambda$  for Game 2 of the Ochs experiment.

- Game 2: both  $A_1$  and  $B_1$  are overplayed (Nash: 0.5 and 0.1)  
 QRE: for intermediate  $\lambda$  precisely this is predicted  
 Note:  $\lambda$  increases over time  $\rightarrow$  suggests learning

### Logit Equilibrium and experiments

- Ochs 95 - Non-zero-sum game:  
52 plays, 4 experience levels of 16 periods

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TABLE X  
DATA AND ESTIMATES FOR OCHS, GAME 3

Period	n	Actual		Predicted		$\lambda$	- $\lambda^*$		
		$A_1$	$B_1$	$\hat{A}_1$	$\hat{B}_1$		QRE	Nash	Rand
1-16	128	0.527	0.366	0.615	0.383	1.856	1747	1822	1774
17-32	128	0.573	0.393	0.610	0.405	1.568	1735	1870	1774
33-48	128	0.610	0.302	0.614	0.301	3.306	1640	1708	1774
48-52	128	0.455	0.285	0.500	0.200	$\infty$	1679	1679	1774
All	512	0.542	0.336	0.619	0.331	2.656	6864	7079	7098

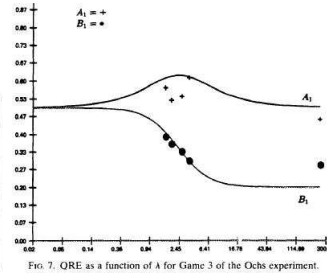


FIG. 7. QRE as a function of  $\lambda$  for Game 3 of the Ochs experiment.

- Game 3: both  $A_1$  and  $B_1$  are overplayed (Nash: 0.5 and 0.2)  
QRE: for intermediate  $\lambda$  precisely this is predicted  
Note:  $\lambda$  increases over time  $\rightarrow$  suggests learning

### Logit equilibrium - Continuous strategies Anderson, Goeree, & Holt, SEJ 2002

- Conceptual issues
- Theoretical results
- Applications:
  - The travelers dilemma
  - Minimal effort game
  - Public goods game

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## Conceptual issues

- As in the original approach of MP95:
  - If decision  $x$  has expected payoff  $\pi_i^e(x)$  then a player  $i$  chooses the decision with the highest value for  $U_i(x) = \pi_i^e(x) + \mu\epsilon_{ix}$ .
  - $\mu$  is a positive ‘error’ parameter and  $\epsilon_{ix}$  is the realization of a random variable
  - Result: choice is stochastic and distribution of the random variable determines the form of the choice probabilities.
  - Double exponential distribution (extreme values)  $\rightarrow$  Logit model

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## Conceptual issues - choice density

- Logit model: choice probabilities are proportional to exponential functions of expected payoffs
- In particular,  $prob_i(x)$  is proportional to  $e^{\pi_i^e(x)/\mu}$   
Note: Higher  $\mu$  makes choice probabilities less sensitive to expected payoff (compare with  $1/\lambda$  in MP95)
- Continuum of choice possibilities  $x \in [\underline{x}, \bar{x}]$  then Logit model specifies the following choice density:

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$$f_i(x) = \frac{e^{\pi_i^e(x)/\mu}}{\int_{\underline{x}}^{\bar{x}} e^{\pi_i^e(y)/\mu} dy}, \quad \forall i \quad (0.1)$$

- Note: this is a ‘continuous’ version of the logistic quantal response function in MP95.

### Conceptual issues - equilibrium

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- Application of choice models to games: take into account that distributions of other's decisions enter the payoff function in choice density function (??)
- Hence, consistency requirement as for Nash equilibrium: belief distributions that determine expected payoffs (r.h.s.) match the decision distributions (l.h.s)
- Thus, the logit choice rule (??) determines the equilibrium distributions as a fixed point  $\rightarrow$  the QRE in Logit form known from MP95.

$$f_i(x) = \frac{e^{\pi_i^e(x)/\mu}}{\int_{\underline{x}}^{\bar{x}} e^{\pi_i^e(y)/\mu} dy}, \quad \forall x \quad \forall i \quad (0.2)$$

### Conceptual issues - equilibrium density

$$f_i(x) = \frac{e^{\pi_i^e(x)/\mu}}{\int_{\underline{x}}^{\bar{x}} e^{\pi_i^e(y)/\mu} dy}, \quad \forall x \quad \forall i$$

Differentiating (both sides of) the logit choice rule (??) w.r.t. the decision  $x$  establishes that the equilibrium densities satisfy:

$$\pi_i^{e'}(x) f_i(x) - \mu f_i'(x) = 0, \quad \forall i \quad (0.3)$$

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the so-called *logit differential equation* in the equilibrium choice density.

- Note: (??) dictates that
  - (i) if  $\mu \rightarrow 0$  then  $\pi_i^{e'}(x) f_i(x) = 0$ ; i.e., the condition for an interior NE: either FOC for payoff maximization is satisfied at  $x$  else  $x$  is chosen with zero probability,  $f_i(x) = 0$ .
  - (ii) if  $\mu \rightarrow \infty$  then noise effect dominates and  $f_i' = 0$ ; i.e., the equilibrium density is uniform.

### Conceptual issues - Rank-based payoffs

- Rank-based payoffs - payoff depends on whether the decision of the player is 'higher' or 'lower' than the decision(s) of the other player(s)

Prominent example: auctions

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- Consider two-person game: let  $\alpha_H(x)$  and  $\alpha_L(x)$  be payoff components associated with player  $i$ 's *own* decision when the decision is higher or lower than the other player's decision.
- Similarly, let  $\beta_H(y)$  and  $\beta_L(y)$  be payoff components associated with the *other* player's decision when the player  $i$ 's own decision is higher or lower ranked.

### Conceptual issues - Rank-based payoffs

- Then the payoff function can be written as

$$\pi_i^e(x) = \int_{\underline{x}}^x [\alpha_H(x) + \beta_H(y)] f_j(y) dy + \int_x^{\bar{x}} [\alpha_L(x) + \beta_L(y)] f_j(y) dy \quad (0.4)$$

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- It can be shown that if the component payoff functions are (i) additively separable and (ii) continuous in own,  $x$ , and other's,  $y$ , decision, then the derivative  $\pi_i^{e'}(x)$  depends only on the player's own decision  $x$ , the other player's distribution and density functions evaluated at  $x$ ,  $F_j(x)$  and  $f_j(x)$ , and possibly some exogenous parameter  $\rho$ .
- Hence, the expected payoff derivative can be written as  $\pi_i^{e'} = \pi_i^{e'}(F_j(x), f_j(x), x, \rho)$ , and satisfies the so-called *local payoff property*.

### Conceptual issues - Rank-based payoffs

- Extension to n-player case - if one's payoff depends on whether one has the highest (or lowest) decision
  - If having highest is critical (as in an auction), then  $H$  and  $L$  in (??) represents the case where one's decision is highest or not
- Density  $f_j(x)$  used in the integrals is replaced by the density of the maximum of the  $n - 1$  other decisions
- Example: Second price auction with values  $v \rightarrow \alpha_H(x) = v, \beta_H(y) = -p_y$ , and  $\alpha_L(x) = \beta_L(y) = 0$ .
- It can be shown that for this case the expected payoff derivative also satisfies the *local payoff property*, i.e.,  $\pi_i^{e'} = \pi_i^{e'}(F_{-i}(x), f_{-i}(x), x, \rho)$ .

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### Properties of the Logit equilibrium

#### **Proposition 1** (EXISTENCE) [AGH SEJ02]

*There exists a logit equilibrium for all n-player games with a continuum of feasible decisions when players' expected payoffs are bounded and continuous in others' distribution functions. Moreover, the equilibrium distribution is differentiable.*

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## Properties of the Logit equilibrium

**Proposition 2** (UNIQUENESS) [AGH SEJ02]

*Any symmetric logit equilibrium for a game satisfying the local payoff property is unique if the expected payoff derivative,*

$$\pi_i^{e'} = \pi_i^{e'}(F_{-i}, f_{-i}, x, \rho) \quad \forall i \text{ is either}$$

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- (a) *strictly decreasing in  $x$ , or*
- (b) *strictly increasing in the common distribution function  $F$ , or*
- (c) *independent of  $x$  and strictly decreasing in the common density function  $f$ , or*
- (d) *a polynomial expression in  $F$ , with no terms involving  $f$  or  $x$ .*

## Properties of the Logit equilibrium

**Proposition 3** (COMPARATIVE STATICS FOR A SYMMETRIC EQUILIBRIUM) [Prop. 4, AGH SEJ02]

*Suppose that the shift parameter  $\rho$  increases the marginal expected payoffs, i.e.  $\partial \pi_i^{e'}(F_{-i}, f_{-i}, x, \rho) / \partial \rho > 0 \quad \forall i$ , for a symmetric game satisfying the local payoff property. Then an increase in  $\rho$  yields stochastically higher logit equilibrium decisions if either*

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- (a)  $\partial \pi_i^{e'} / \partial x \leq 0$ , or
- (b)  $\partial \pi_i^{e'} / \partial F_{-i} \geq 0$ .

## An experiment!

### The Travelers' Dilemma

Basu, AER 1994

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#### Logit equilibrium - Travelers' dilemma

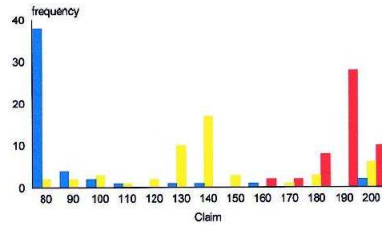
- The Travelers' dilemma (in continuous strategies)
  - Payoffs: two travelers; each is reimbursed the minimum of the claims in  $[\underline{x}, \bar{x}]$  and a fixed penalty  $R > 0$  is transferred from the high claimant to the low claimant
  - The penalty gives each traveler an incentive to 'undercut'  
 $\Rightarrow$  both claiming  $\underline{x}$  is the **unique Nash equilibrium**  
**Independent** of size of penalty  $R$ .

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## Logit equilibrium - Travelers' dilemma

- Travelers' dilemma - the data

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blue: R=50, yellow:R=25, red: R=10;  
frequency of decisions in final 5 of 10  
rounds;  
source: GH PNAS99

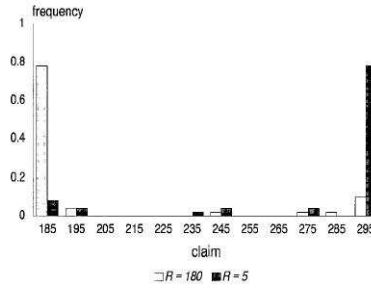


FIGURE 1. CLAIM FREQUENCIES IN A TRAVELER'S DILEMMA FOR  $R = 180$  (LIGHT BARS) AND  $R = 5$  (DARK BARS)

One-shot decisions;  
source: GH AER01

## Logit equilibrium - Travelers' dilemma

$$\pi_i^e(x) = \int_{\underline{x}}^x [\alpha_H(x) + \beta_H(y)] f_j(y) dy + \int_x^{\bar{x}} [\alpha_L(x) + \beta_L(y)] f_j(y) dy$$

- Travelers' dilemma - Theoretical result

Expected payoff has rank-based property:

Let  $x$  be player  $i$ 's own decision, then

$$\alpha_H(x) = -R, \alpha_L(x) = x, \beta_H(y) = y, \beta_L(y) = +R$$

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Hence, the expected payoff of player  $i$  is given by:

$$\pi_i^e(x) = \int_{\underline{x}}^x [-R+y] f_j(y) dy + \int_x^{\bar{x}} [x+R] f_j(y) dy \quad i, j = 1, 2, j \neq i$$

And its derivative is:

$$\pi_i^{e'}(x) = 1 - F_j(x) - 2Rf_j(x) \quad i, j = 1, 2, j \neq i$$

### Logit equilibrium - Travelers' dilemma

$$\pi_i^{e'}(x) = 1 - F_j(x) - 2Rf_j(x) \quad i, j = 1, 2, \quad j \neq i$$

#### Slide 37

- Travelers' dilemma - Interpretation of expected payoff derivative
  - $1 - F_j(x)$  ... probability that the other's claim is higher  $\rightarrow$  unilateral increase of claim raises minimum and, hence, the payoff.
  - $-2R$  ... payoff reduction when  $i$ 's claims crosses over  $j$ 's claim.
  - $f_j(x)$  ... probability with which this cross-over occurs.

### Logit equilibrium - Travelers' dilemma

$$\pi_i^{e'}(F_j, f_j, x, -R) = 1 - F_j(x) - 2Rf_j(x) \quad i, j = 1, 2, \quad j \neq i$$

#### Slide 38

- Travelers' dilemma - Property of equilibrium
  - Expected payoff derivative has local payoff property.
  - Expected payoff derivative increases with parameter  $\rho := -R$ .
  - Expected payoff derivative does not increase with  $x$ :  $\partial \pi_i^{e'} / \partial x = 0$ .
  - Hence, Proposition ?? implies that a decrease in  $R$  yields (stochastically) higher claims.

This is **precisely what we observe!** And, is **not predicted** by the standard **Nash equilibrium**.

### Logit equilibrium - Travelers' dilemma

$$\pi_i^{el}(x)f_i(x) - \mu f_i'(x) = 0, \forall i$$

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- Travelers' dilemma - Numerical solution
- Substituting the expected payoff derivative into the logit differential equation (see box) gives  $\rightarrow$  second order differential equation in the equilibrium distribution  $F(x)$ :

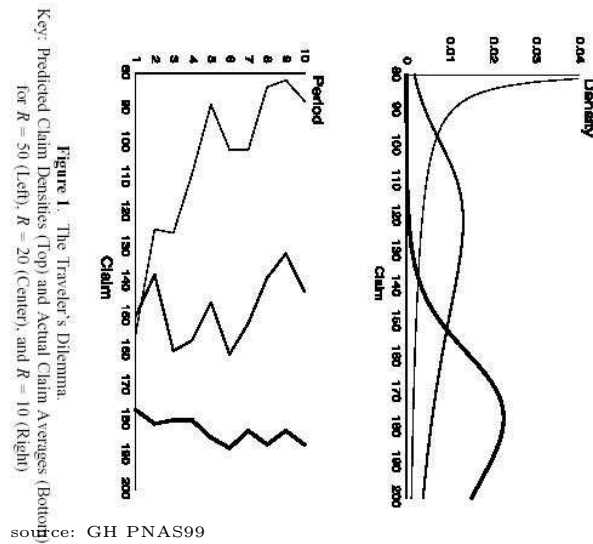
$$[1 - F_j(x) - 2Rf_j(x)]f_i(x) - \mu f_i'(x) = 0, \forall i, j = 1, 2, i \neq j$$

There exists no analytical solution but it can be solved numerically for given noise  $\mu$ .

### Logit equilibrium - Travelers' dilemma

- Travelers' dilemma - Numerical solutions  $\mu = 8.5$  and actual

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### Logit equilibrium - Coordination game

- Minimum-effort coordination game
  - Payoff of a player: minimum effort of all players minus costs  $c$ ,  $0 < c < 1$ , of the player's own effort,  $x \in [0, \bar{x}]$ .  
Player  $i$ 's expected payoff:

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$$\pi_i^e(x) = \int_0^x y f_j(y) dy + \int_x^{\bar{x}} x f_j(y) dy - cx, \quad \forall i, j = 1, 2, i \neq j$$

First part: own effort  $x$  is above other's effort; second part: own effort  $x$  is minimum effort.

Expected payoff derivative:

$$\pi_i^{e'}(F_j, f_j, x, -c) = 1 - F_j(x) - c, \quad i, j = 1, 2, j \neq i$$

### Logit equilibrium - Coordination game

$$\pi_i^{e'}(F_j, f_j, x, -c) = 1 - F_j(x) - c, \quad i, j = 1, 2, j \neq i$$

- Minimum-effort coordination game - Interpretation of expected payoff derivative
    - $1 - F_j(x)$  ... probability that other's effort is higher  $\rightarrow$  unilateral increase of effort raises minimum effort and, hence, the payoff.
    - $-c$  ... cost of raising effort.
  - Minimum-effort coordination game - Property of equilibrium
    - $\pi_i^{e'} > 0 (< 0)$  if  $F_j(x) = 0 (= 1) \Rightarrow$  all effort levels are NE.
    - Expected payoff derivative increases with parameter  $\rho := -c$ .
    - Expected payoff derivative does not increase with  $x$ .
- Hence, Proposition ?? implies that a decrease in effort costs  $c$  yields (stochastically) higher effort.

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This is **intuitive** but **not predicted** in **Nash equilibrium**.

### Logit equilibrium - Coordination game

- Minimum-effort coordination game - some data

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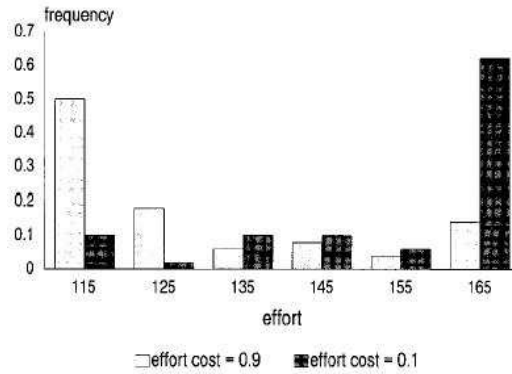


FIGURE 2. EFFORT CHOICE FREQUENCIES FOR A MINIMUM-EFFORT COORDINATION GAME WITH HIGH EFFORT COST (LIGHT BARS) AND LOW EFFORT COST (DARK BARS)

Source: GH AER01

### Logit equilibrium - Public good game

$$\pi_i^e(x) = \int_{\underline{x}}^x [\alpha_H(x) + \beta_H(y)] f_j(y) dy + \int_x^{\bar{x}} [\alpha_L(x) + \beta_L(y)] f_j(y) dy$$

- Public good game
  - Payoff of a player depends on own contribution  $x_i$  and on the sum of the other's contribution  $\sum_{j \neq i} x_j$ :

Player  $i$ 's payoff:

$$\pi_i(x_i) = w_i - x_i + R_I x_i + R_E \sum_{j \neq i} x_j$$

$R_I < 1$  ... internal return from own contribution,  
 $R_E$  ... external return from the sum of other's contribution

Note: This payoff is a trivial special case of the rank-based payoff; is independent from whether or not one's own contribution is highest or not.

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## Logit equilibrium - Public good game

$$\pi_i(x_i) = w_i - x_i + R_I x_i + R_E \sum_{j \neq i} x_j$$

- Public good game - Properties of equilibrium

Marginal expected payoff from own contribution is constant:  $R_I - 1$ : Hence,

Proposition ?? (d) implies **uniqueness**

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The marginal expected payoff is increasing in the exogenous parameter  $R_I$  and non-increasing in own contribution  $x$ : Hence,

Proposition ?? (a) implies a stochastic **increase in contributions** with an **increase in  $R_I$** .

This is what we **see in experiments** but is **not predicted** in **Nash equilibrium**.

## Combining Inequity aversion and Noise

Goeree & Holt, EER 2000

- Alternating offer bargaining
  - Two-stage alternating offer bargaining game

Initial pie size  $s = 2.40$

Treatments: different remaining pie sizes after rejection  
different fixed payments to proposer and responder

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Parameters for the 7 bargaining games							
Remaining pie size (\$R)	0.00	0.40	0.80	1.20	1.60	2.00	2.40
Fixed payment to the initial proposer	2.65	2.25	1.85	1.45	1.05	0.65	0.25
Fixed payment to the initial responder	0.25	0.65	1.05	1.45	1.85	2.25	2.65

Standard SPE: proposer offers  $R$ , which is accepted  
However, fixed payments **exaggerate earnings inequalities** in equilibrium

## Combining Inequity aversion and Noise

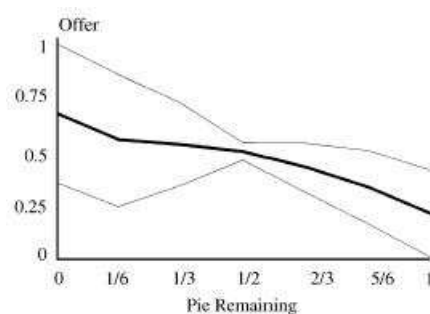
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- Alternating offer bargaining - Predictions
  - Subgame perfect Nash equilibrium → perfect positive correlation between initial offers and  $R$
  - Final earnings equalization → perfect negative correlation between initial offers and  $R$
  - Note: These relationships might be flattened by noise, but any negative correlation indicates fairness considerations
  - Procedures: 6 sessions, each proposer made 7 proposals, 1 for each treatment in random order, sequentially; then random assignment of proposals to responders, sequentially; accepted or rejected and made counter offer; matched proposer accepted or rejected; all divisions paid out

## Combining Inequity aversion and Noise

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- Alternating offer bargaining - Results



Average 1st stage offers with std.dev.;  
Source: GH EER00

Clear negative relation between remaining pie size and offers. Standard prediction of positive correlation clearly rejected.

Note 1: Rabin's intention model also predicts positive correlation.

Note 2: The pure Logit model can also not explain the negative correlation.

### Combining Inequity aversion and Noise

- Alternating offer bargaining - FS inequity aversion with noise
  - FS inequity aversion without noise:

For illustration consider

$$R = 0, FP_R = 2.65, FP_R = 0.25 \Rightarrow \pi_P > \pi_R \quad \forall \text{ offers } x:$$

$$U_P(x) = 2.65 + 2.40 - x - \beta_P[2.40 - 2x]$$

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If  $\beta > 1/2$ :  $U'_P(x) > 0 \Rightarrow x^* = 2.40$  (the whole pie)

If  $\beta < 1/2$ :  $U'_P(x) < 0 \Rightarrow x^*$  depends on  $\alpha_R$  but is always less than half of the pie,  $2.40/2 = 1.20$ .

Actual data: three offered whole pie, six offered less than half of the pie, but 13 (out of 23) made offers outside predicted range.

### Combining Inequity aversion and Noise

- Alternating offer bargaining - FS inequity aversion with noise
  - FS inequity aversion with noise:

Use logit decision model but replace expected material utility with FS utility:

Assume all responders have same  $\alpha_R$ ,  $\beta_R$ , and  $\mu_R$

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Assume all proposers have same  $\alpha_P$ ,  $\beta_P$ , and  $\mu_P$

Estimate these parameters from the data with maximum likelihood technique, to get

Maximum-likelihood estimates

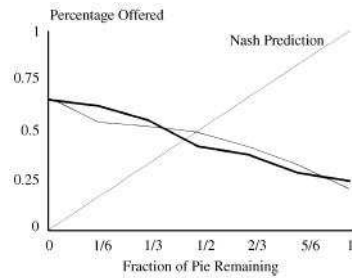
Variable	$\mu_P = \mu_R$	$\alpha_P = \alpha_R$	$\beta_P$	$\beta_R$
Estimate	0.55	0.84	0.66	0.12
(Standard error)	(0.06)	(0.16)	(0.08)	(0.02)

Source: GH EER00

### Combining Inequity aversion and Noise

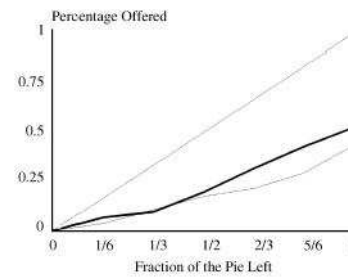
- Alternating offer bargaining - FS inequity aversion with noise

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Initial offers;

Source: GH EER00



Counter offers

Source: GH EER00

Thick lines: prediction; thin lines: actual averages

Hence, the model fits the data quite well.